

# search

## Agents

Agents interest of  
Complete Process  
(simple system)

Agents interest of what  
will happen after ending  
the Process of Find  
best solution.  
(complex system)

### 1) Reflex agent (simple system)

- a) choose action based on current Percept.
- b) may have memory or model of world's current state.
- c) Do not consider future consequence of their actions.  
"Do not ask ~~if~~ what if"
- d) Consider how the world is (cannot care about future actions)
- e) Can it be rational (yes)?  
→ yes, if it is provided by if conditions.

## 2] Planning Agents (Complex System)

- a) Ask "what-if"
- b) Decisions based on (hypothesized) consequences of actions.
- c) must have model of how world evolves in response to actions.
- d) must Formulate a goal (test)  $\Rightarrow$  must reach the goal.
- e) Consider how the world would be.

### Planning

- $\rightarrow$  Agent make Plan according to search.
- $\rightarrow$  It take period of time to Plan then take decision according to Plan.
- $\rightarrow$  Plan only one time

### RePlanning

- $\rightarrow$  Agents make Plan for the first step according to search then take decision to complete first step. then
- $\rightarrow$  It replan for next step according to new search and take new best decision to start work for next step.
- $\rightarrow$  Plan for each step  
 $\hookrightarrow$  no. of plans = no. of steps.



## \* optimal vs Complete Planning

↳ Agent, not only reach goal or take action but also it want to get optimal solution of problem. (search for best way to goal or minimum time to reach goal)

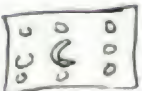
## \* Search Problems:

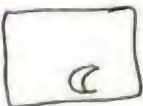
↳ Process to take right decision.

It consists of

1) state space (world): All possible cases of the model.

ex mouse want to eat 9 piece of cheese.

 → mouse eat piece on center

 → eat the last piece.

2) A Successor Function (actions, costs)

↳ For any state what action can I take and what cost of taking it.

### 3) A start state and goal test

start state: state from which i can start working.

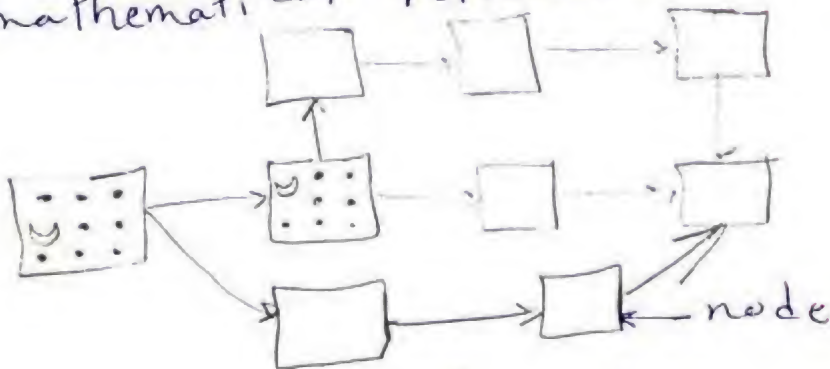
goal test: the final state.

→ solution of search problem is sequence of action which transforms start state to goal state.

### \* Uniformed search methods

#### ① state space <sup>graph</sup> representation

→ mathematical representation of search problem.



- nodes → world configurations.
- Arc → successors (action results)
- goal test is set of goal nodes.

→ here we represent needed states only.

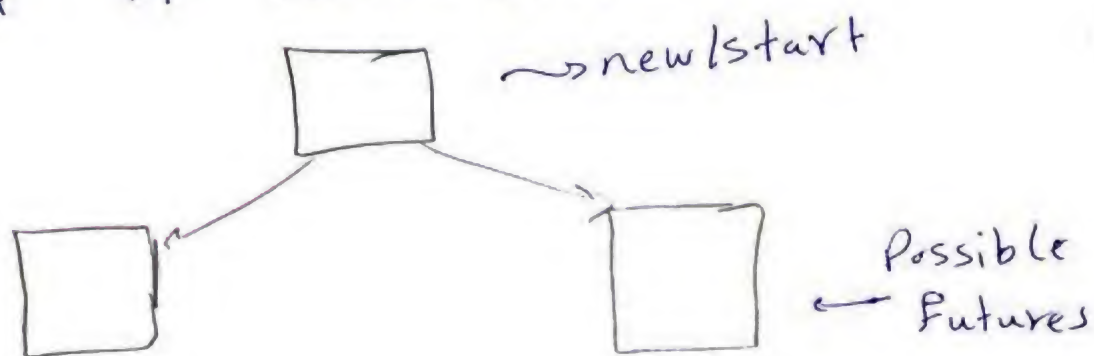
→ we can rarely build this full graph.

→ there is no start state. but goal test.



## ② Search trees

↳ part of search graph.



### Search tree

- start state is the root node.
- children correspond to successors.
- we can never actually build whole tree.

### x Searching with search tree

- try to expand as few tree nodes as possible.
- maintain a Fringe of partial plans under consideration.
- try to expand as few tree nodes as possible

(General Tree Search)

Fringe : all of plans that may yet work.

→ all ways that can reach me from start to goal during world state search tree.

→ one of this Frings is my solution (optimal solution)

### Expansion

↳ Picking some thing out of the Fringe and if it is not a goal already.

→ Process of expanding new state from current state on one of Frings. we can't do this process if we reach goal.

### \* Exploration strategy

↳ what Fringe nodes do you explore next?  
↳ what will happen next?

### \* Main Question:-

↳ which Fringe nodes to explore to reach goal?

### Depth-First Search

strategy: expand deepest node first

Implementation: Fringe is LIFO stack.

Solution: left most solution.



## Search Algorithm Properties

→ For any Algorithm we check for 4 Properties:-

[1] Complete:

↳ Ability to reach goal if it exists.

[2] Optimal:

↳ Ability to reach to best solution.

[3] Time Complexity

↳ time to reach to solution  
↳ time to expand all nodes of one Fringe equal  
to time to expand one node of fringe  $\times$  no. of  
Fringe nodes.

[4] Space Complexity

↳ size of Fringe in memory. size of  
nodes of Fringe in stack.

\* no. of nodes in entire tree.

$$1 + b + b^2 + \dots + b^m = O(b^m)$$

↳ branching factor: start node can make  
a branch for  $b$  number of children.

# Depth First Properties

## [1] Time Complexity

- what nodes DFS expand?
  - ↳ some left prefix of tree.
  - ↳ could process the whole tree!
    - ↳ worst case expand all nodes (right most)
    - ↳ best case expand no nodes (start  $\equiv$  goal)

worst case:  $O(b^m)$  , best case  $\approx O(1)$

## [2] Space Complexity

- ↳ only has siblings on path to root, so  $O(b^m)$
- ↳ worst case we get solution on last tier.

## [3] Complete

- ↳ m could be finite, if we prevent cycles.
- solution can only be found if
  1. it exists
  2. finite algorithm.

## [4] optimal

- no, it finds left most solution regardless of depth or cost.



## [2] Breadth - First (BFS)

Strategy: expand a shallowest node first.

Implementation: fringe is FIFO queue.

Solution: shallowest solution (shortest fringe)

### \* \*DFS & BFS

→ BFS will outperform DFS when:

- 1) need less time complexity.
- 2) need shallowest solution.
- 3) need optimal solutions for costs equal to 1.
- 4) need complete solutions.

→ DFS will outperform BFS when:

- 1) most solutions at left side of tree.
- 2) all solutions at the last level.
- 3) need less space complexity.

[3]

### [3] uniform cost

strategy: expand a cheapest node first.

Implementation: Fringe is priority queue

Solution: cheapest solution.

→ expand node of cheapest cost.

#### \* Advantages

→ complete and optimal.

#### \* Disadvantages

1. no information about goal location.

(expand)  $\text{cost}(n)$  (Goal)  $\text{cost}(n) \sim \text{cost}(n) \sim \text{cost}(n)$   
•  $\text{cost}(n) \sim \text{cost}(n)$

2. explores options in every direction.

↳ no determined direction on our work.

(10)



## \* The one Queue: Priority queues

- Conceptually, all frings are priority queues.
- For DFS and BFS you can avoid the  $\log(n)$  overhead from actual priority queue with stacks and queues.
- Can even code one implementation that takes variable queuing object.

## Informed search

→ we need to know information about goal location to know if I work on correct way or not.

→ we have one function, 2 search Algorithms:

### 1) Heuristics

→ function takes state of state space, and give number which represent how far the goal location from my location to doing process.

## \* Greedy search

- ↳ Search Algorithm use idea of Heuristic.
- ↳ It is not an optimal search Algorithm.

## \* A\* Algorithm search

- ↳ It collects all last search Algorithms ideas to get very good search Algorithm.

## \* Graph search

### (Recap search)

#### 1) Search Problem

- a) states (configuration of world)
- b) actions & costs.
- c) successor Function (say how states respond to actions)
- d) start state & goal state.

#### 2) search tree

- a) nodes: Plans for reaching states.
- b) Plans have costs (sum of action costs)


#### 3) search Algorithm


- a) systematically builds search tree.
- b) chooses an ordering of the fringe.
- c) optimal: find the least-cost Plans.



## EX: Pancake Problem

Problem: need to arrange pancake from big to small.  
states: state of pancakes during flipping to reach goal.  
Costs: no. of pancakes flipped.

start state: 

goal state: 

Algorithm: we can use UCS algorithm according to cost also we can use DFS or BFS.

## Search Heuristics

heuristic is

- Function that estimates how close a state is to goal.
- Designed for particular search problem.
- we make it every step to see ~~how~~ <sup>if</sup> close to goal or not.
- every search problem need different heuristic according to natural of problem.

$\therefore$  no. of largest pancake

Heuristic: no. of largest pancake that is still out of place.

## Greedy search

strategy: expand the nodes that seems closest to goal according to Heuristic.

Heuristic: estimate the distance to nearest goal for each state.

Implementation: Pring is Priority queue &:

solution: best heuristic solution.

→ It doesn't care about cost.

→ cares only about heuristic

~~Common~~ case: Best-First takes you to wrong goal.

worst case: like badly-guided DFS.

solution → Combine UCS and Greedy search.

## A\* search

→ Combining UCS and Greedy Search.

\* uniform cost orders by path cost or backward cost  $g(n)$ .

\* Greedy search " by goal proximity or forward cost  $h(n)$ .



→ should we stop when we enqueue goal?  
→ No only stop when we dequeue goal.

### Admissibility

- \* Inadmissible (pessimistic) heuristics break optimality by trapping good plans on fringe.
- \* Admissible (optimistic) heuristics slow down bad plans but plans never outweigh true costs.

### (Admissible heuristic)

→ A heuristic  $h$  is admissible if:

$$0 \leq h(n) \leq h^*(n) \rightarrow h^*(n) = g(n)$$

→  $h^*(n)$  is true cost to nearest goal

$h(n) > h^*(n)$  → (optimal solution) → مستبعد، نوبت

### \* A\* applications

- \* video games.
- \* language analysis.
- \* Machine translation
- \* speech recognition.
- \* Robot motion Planning.
- \* resource planning problems.

## \* Creating Admissible heuristics

→ most of work in solving hard search Problems optimally is in coming up with admissible heuristics.

→ often, admissible heuristics are solutions to relaxed problems where new actions are available.

→ inadmissible heuristics are often useful too.

EX 8 Puzzle

7	2	4
5		6
8	3	1

start state

3	7	1
2	4	5
		6

Action

	1	2
3	4	5
6	7	8

Goal state

states? all cases that i can move any number from its location to another one if it is next to free space

How many states →  $8! =$

actions? moving number to free space next to it.



\* no of successors from start state?  
maximum 4 moves (minimum 2 moves)

\* Cost should be? 1 every time

\* Heuristic: no. of tiles misplaced.

$$h(\text{start}) = 8$$

\* why it is admissible?

1) There is relaxed - problem heuristic  
(easiest solution but need more work)

2) Direct move  
↳ move every number from start state to final state.

$h = \text{no. of moves.}$

$h(1) = 3$  (moves to reach to its correct state at goal state)

$h(2) = 1$  ,  $h(3) = 2$  and so on.

How Total  $h = 18$  From start to goal.

Admissible if  $0 \leq h(n) \leq 2(n)$

→ How about using actual cost as heuristic?

$$f(n) = h(n) + g(n) = 2g(n) = 2h(n)$$

if  $g(n) = h(n)$

→ we go back to UCS by double cost  
value trade off: complex work, neglect heuristic  
with  $A^*$  → trade off between quality of  
estimate and work per-node.

### \* Trivial heuristic

• bottom of lattice is zero heuristic

• top " " is the exact "

if  $(h=0)$  no heuristic → go back to  $\left( \begin{array}{l} \text{uniform} \\ \text{cost} \\ \text{search} \end{array} \right)$

Dominance:  $h_a \leq h_c$

$\forall n: h_a(n) \geq h_c(n)$  For all nodes

we choose  $h_a(n)$  (best one)

→ Heuristics From a semi-lattice

→ max of admissible heuristics is admissible

$$h(n) = \max(h_a(n), h_b(n))$$

we choose



## \* Graph search

↳ Failure to detect repeated states can cause exponential more work.

idea: never expand state twice.

Implementation:

- \* tree search + set of expanded states (closest set)
- \* Expand search tree node by node, but never expand node twice.
- \* Before expanding a node, check if it is expanded before or not.
- \* if not now, skip it, if new add to closest set.

Important

- \* store closest set as a set, not a list.
- \* complete  $\rightarrow$  if goal exists, we will find it.
- \* optimal  $\rightarrow$  no.

(19)

	Depth First search	Breadth First search	Uniform Cost search
nodes expanded	<ul style="list-style-type: none"> <li>→ some left prefix of the tree.</li> <li>→ Could process the whole tree.</li> <li>→ if <math>m</math> is finite, take time <math>O(b^m)</math></li> </ul>	<ul style="list-style-type: none"> <li>→ processes all nodes above shallowest solution.</li> <li>→ let depth of shallowest solution be <math>s</math>.</li> <li>→ Search takes time <math>O(b^s)</math></li> </ul>	<ul style="list-style-type: none"> <li>→ process all nodes with cost less than cheapest solution.</li> <li>→ takes time <math>O(\frac{C^*}{\epsilon})</math></li> <li><math>C^* \rightarrow</math> sol. costs</li> <li><math>\epsilon</math>.</li> </ul>
space <del>done</del> that brings take	$O(b^m)$	$O(b^s)$	$O(\frac{C^*}{\epsilon})$
complete	only if we prevent cycles	$s \rightarrow$ finite Yes	Assume best sol. has finite cost and min. $\epsilon$ $\rightarrow$ Yes
optimal	No	only if costs are all 1	Yes



→ heuristic

- \* Function estimates how close a state is to a goal.
- \* Designed for Particular search Problem.
- \* EX: Manhattan distance.

### Greedy search

\* strategy expand a node that you think is closest to a goal state.

↳ heuristic  $\Rightarrow$  estimate of distance to nearest goal for each state.

\* Common case

↳ Best-first<sup>+</sup> takes you to wrong goal.

### A\* search

↳ only stop when we dequeue a goal.

→ Admissible Heuristics

→ heuristic  $h$  is admissible (optimistic) if

$$0 \leq h(n) \leq h^*(n)$$

Where  $\Rightarrow h^*(n)$  is true cost to nearest goal.

## UCS vs $A^*$ Contours

UCS  $\rightarrow$  expands equally in all "directions".

$A^*$   $\rightarrow$  expands mainly toward the goal but does hedge its bets to ensure optimality

## $A^*$ Applications

Video games, language analysis, speech ~~recognition~~ recognition, Machine translation, Pathing.

$\rightarrow$  search problem consists of

1) state space

2) successor function  
(with actions, costs)

3) start state and goal test.

solution  $\rightarrow$  is sequence of actions (a plan)  
which transform start state to goal state.